

Note

The maximum degree of a minimally hamiltonian-connected graph



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ABSTRACT

We determine the possible maximum degrees of a minimally hamiltonian-connected graph with a given order. This answers a question posed by Modalleliyan and Omoomi in 2016. We also pose two unsolved problems.

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We consider finite simple graphs and follow the book [5] for terminology and notations. The *order* of a graph is its number of vertices, and the *size* is its number of edges. We denote by $V(G)$ and $E(G)$ the vertex set and edge set of a graph G respectively. For two graphs G and H , $G \vee H$ denotes the *join* of G and H , which is obtained from the disjoint union $G + H$ by adding edges joining every vertex of G to every vertex of H . K_n and C_n denote the complete graph of order n and the cycle of order n respectively. The *wheel* of order n , denoted by W_n , is the graph $K_1 \vee C_{n-1}$.

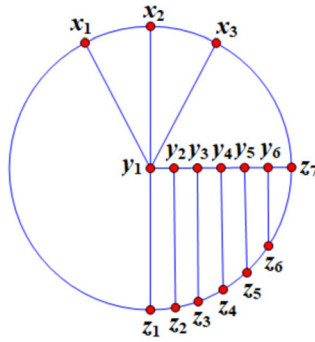
A graph is called *hamiltonian-connected* if between any two distinct vertices there is a Hamilton path. Obviously, any hamiltonian-connected graph of order at least 4 is 3-connected, and hence has minimum degree at least 3.

Definition. A hamiltonian-connected graph G is said to be *minimally hamiltonian-connected* if for every edge $e \in E(G)$, the graph $G - e$ is not hamiltonian-connected.

Clearly every hamiltonian-connected graph contains a minimally hamiltonian-connected spanning subgraph. Concerning the maximum degree of a minimally hamiltonian-connected graph with a given order, Modalleliyan and Omoomi [2] proved the following results: (1) The maximum degree of any minimally hamiltonian-connected graph of order n is not equal to $n - 2$; (2) the wheel W_n is the only minimally hamiltonian-connected graph of order n with maximum degree $n - 1$; (3) for every integer $n \geq 6$ and any integer Δ with $\lceil n/2 \rceil \leq \Delta \leq n - 3$, there exists a minimally hamiltonian-connected graph of order n with maximum degree Δ . They [2] posed the question of whether for Δ in the range $3 \leq \Delta < \lceil n/2 \rceil$, there exists a minimally hamiltonian-connected graph of order n with maximum degree Δ . In this note we answer the question affirmatively. Our construction covers the whole range $3 \leq \Delta \leq n - 1$, not only $3 \leq \Delta < \lceil n/2 \rceil$.

Theorem 1. Let $n \geq 4$ be an integer. There exists a minimally hamiltonian-connected graph of order n with maximum degree Δ if and only if $3 \leq \Delta \leq n - 1$ and $\Delta \neq n - 2$, where $\Delta = 3$ occurs only if n is even.

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Fig. 1. The graph $G(16,5)$.

Proof. Suppose there exists a minimally hamiltonian-connected graph G of order n with maximum degree Δ . Then G is 3-connected, implying that $\Delta \geq \delta \geq 3$ where δ denotes the minimum degree of G . Modalleliyan and Omoomi [2] proved that $\Delta \neq n - 2$. If $\Delta = 3$, then G is cubic and hence its order n is even.

Conversely suppose $3 \leq \Delta \leq n - 1$ and $\Delta \neq n - 2$, and when $\Delta = 3$, n is even. We will construct a minimally hamiltonian-connected graph of order n with maximum degree Δ . If $\Delta = n - 1$, the wheel graph W_n is a minimally hamiltonian-connected graph of order n with maximum degree $n - 1$. Next suppose $3 \leq \Delta \leq n - 3$. We will distinguish the two cases when $n - \Delta$ is odd and when $n - \Delta$ is even. For symbols such as x_i below, if i exceeds its valid range, then x_i does not appear.

Case 1. $n - \Delta$ is odd

We define a graph $G(n, \Delta)$ as follows. Denote $k = \Delta - 2$ and $s = (n - \Delta + 1)/2$. We have $k \geq 1$ and $s \geq 2$.

$$\begin{aligned} V(G(n, \Delta)) &= \{x_1, x_2, \dots, x_k\} \cup \{y_1, y_2, \dots, y_s\} \cup \{z_1, z_2, \dots, z_{s+1}\}, \\ E(G(n, \Delta)) &= \{x_i x_{i+1} \mid i = 1, \dots, k-1\} \cup \{y_i y_{i+1} \mid i = 1, \dots, s-1\} \cup \{z_i z_{i+1} \mid i = 1, \dots, s\} \\ &\quad \cup \{y_1 x_i \mid i = 1, \dots, k\} \cup \{y_i z_i \mid i = 1, \dots, s\} \cup \{x_1 z_1, x_k z_{s+1}, y_s z_{s+1}\}. \end{aligned}$$

The graph $G(16, 5)$ is depicted in Fig. 1.

Clearly the graph $G(n, \Delta)$ has order n and maximum degree Δ . We first show that the graph $G(n, \Delta)$ is hamiltonian-connected; i.e., for any two distinct vertices u and v , there is a Hamilton (u, v) -path. There are 8 cases for the vertex pairs (u, v) . In each case we display a Hamilton (u, v) -path.

We define several symbols to describe the Hamilton paths. For i and j with $i \leq j$, $x_i \overrightarrow{X} x_j$, $y_i \overrightarrow{Y} y_j$ and $z_i \overrightarrow{Z} z_j$ denote the paths $x_i x_{i+1} \dots x_j$, $y_i y_{i+1} \dots y_j$ and $z_i z_{i+1} \dots z_j$, respectively. Throughout $h \overrightarrow{P} g$ will denote a path starting from the vertex h and ending at the vertex g whose vertex set is to be specified. For $1 \leq i \leq s$, let $x_1 \overrightarrow{P} y_i$ and $x_1 \overrightarrow{P} z_i$ denote the paths both with vertex set $\{x_1, y_1, \dots, y_i, z_1, \dots, z_i\}$. Similarly, let $y_i \overrightarrow{P} z_{s+1}$ and $z_i \overrightarrow{P} z_{s+1}$ denote the paths both with vertex set $\{y_i, \dots, y_s, z_i, \dots, z_s, z_{s+1}\}$. Also, we let $z_{s+1} \overrightarrow{P} z_{s+1} = z_{s+1}$. Finally, given a sequence $p \overrightarrow{S} q$, we denote by $q \overleftarrow{S} p$ its reverse sequence. In the following Hamilton paths, strings like $a_i \overrightarrow{A} a_j$ or $a_j \overleftarrow{A} a_i$ with $i > j$ do not appear.

Case 1.1. (x_i, x_j) with $1 \leq i < j \leq k$.

$$x_i \overrightarrow{X} x_{j-1} y_1 x_{i-1} \overleftarrow{X} x_1 z_1 z_2 \overrightarrow{P} z_{s+1} x_k \overleftarrow{X} x_j.$$

Case 1.2. (x_i, y_j) with $1 \leq i \leq k$ and $1 \leq j \leq s$.

$$x_i \overrightarrow{X} x_k z_{s+1} \overleftarrow{P} z_{j+1} z_j \overleftarrow{Z} z_1 x_1 \overrightarrow{X} x_{i-1} y_1 \overrightarrow{Y} y_j.$$

Case 1.3. (x_i, z_j) with $1 \leq i \leq k$ and $1 \leq j \leq s$.

$$x_i \overrightarrow{X} x_k z_{s+1} \overleftarrow{P} y_{j+1} y_j \overleftarrow{Y} y_1 x_{i-1} \overleftarrow{X} x_1 z_1 \overrightarrow{Z} z_j,$$

where when $j = s$ the string $z_{s+1} \overleftarrow{P} y_{j+1}$ means z_{s+1} .

Case 1.4. (x_i, z_{s+1}) with $1 \leq i \leq k$.

$$x_i \overrightarrow{X} x_k y_1 x_{i-1} \overleftarrow{X} x_1 z_1 z_2 \overrightarrow{P} z_{s+1}.$$

Case 1.5. (y_i, y_j) with $1 \leq i < j \leq s$.

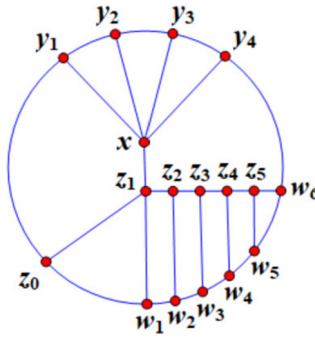
$$y_i \overrightarrow{Y} y_{j-1} z_{j-1} \overleftarrow{Z} z_i z_{i-1} \overleftarrow{P} x_1 x_2 \overrightarrow{X} x_k z_{s+1} \overleftarrow{P} y_j,$$

where when $i = 1$ the string $z_{i-1} \overleftarrow{P} x_1$ means x_1 .

Case 1.6. (y_i, z_j) with $1 \leq i < j \leq s + 1$.

$$y_i \overrightarrow{Y} y_{j-1} z_{j-1} \overleftarrow{Z} z_i z_{i-1} \overleftarrow{P} x_1 x_2 \overrightarrow{X} x_k z_{s+1} \overleftarrow{P} z_j,$$

where when $i = 1$ the string $z_{i-1} \overleftarrow{P} x_1$ means x_1 .

Fig. 2. The graph $H(17,5)$.

Case 1.7. (y_i, z_j) with $1 \leq j \leq i \leq s$.

$$y_i \xrightarrow{\overleftarrow{Y}} y_j y_{j-1} \xrightarrow{\overleftarrow{P}} x_1 x_2 \xrightarrow{\overrightarrow{X}} x_k z_{s+1} \xrightarrow{\overleftarrow{P}} z_{i+1} z_i \xrightarrow{\overleftarrow{Z}} z_j,$$

where when $j = 1$ the string $y_{j-1} \xrightarrow{\overleftarrow{P}} x_1$ means x_1 .

Case 1.8. (z_i, z_j) with $1 \leq i < j \leq s + 1$.

$$z_i \xrightarrow{\overrightarrow{Z}} z_{j-1} y_{j-1} \xrightarrow{\overleftarrow{Y}} y_i y_{i-1} \xrightarrow{\overleftarrow{P}} x_1 x_2 \xrightarrow{\overrightarrow{X}} x_k z_{s+1} \xrightarrow{\overleftarrow{P}} z_j,$$

where when $i = 1$ the string $y_{i-1} \xrightarrow{\overleftarrow{P}} x_1$ means x_1 .

We have shown that $G(n, \Delta)$ is hamiltonian-connected. Recall that any hamiltonian-connected graph of order at least 4 is 3-connected and hence has minimum degree at least 3. Note that every edge of $G(n, \Delta)$ has one endpoint of degree 3. Thus for any $e \in E(G(n, \Delta))$, $G(n, \Delta) - e$ has a vertex of degree 2, implying that it is not hamiltonian-connected. This completes the proof that $G(n, \Delta)$ is minimally hamiltonian-connected.

Case 2. $n - \Delta$ is even

We define a graph $H(n, \Delta)$ as follows. Denote $k = \Delta - 1$ and $s = (n - \Delta - 2)/2$. Since $n - \Delta$ is even and $\Delta \leq n - 3$, we have $k \geq 3$ and $s \geq 1$.

$$\begin{aligned} V(H(n, \Delta)) &= \{x\} \cup \{y_1, y_2, \dots, y_k\} \cup \{z_0, z_1, z_2, \dots, z_s\} \cup \{w_1, w_2, \dots, w_{s+1}\}, \\ E(H(n, \Delta)) &= \{y_i y_{i+1} \mid i = 1, \dots, k-1\} \cup \{z_i z_{i+1} \mid i = 0, 1, \dots, s-1\} \cup \{w_i w_{i+1} \mid i = 1, \dots, s\} \\ &\quad \cup \{xy_i \mid i = 1, \dots, k\} \cup \{z_i w_i \mid i = 1, \dots, s\} \cup \{xz_1, y_1 z_0, z_0 w_1, y_k w_{s+1}, z_s w_{s+1}\}. \end{aligned}$$

The graph $H(17, 5)$ is depicted in Fig. 2.

Clearly the graph $H(n, \Delta)$ has order n and maximum degree Δ . We first show that the graph $H(n, \Delta)$ is hamiltonian-connected; i.e., for any two distinct vertices u and v , there is a Hamilton (u, v) -path. There are 18 cases for the vertex pairs (u, v) . In each case we display a Hamilton (u, v) -path.

We define several symbols to describe the Hamilton paths. For i and j with $i \leq j$, $y_i \xrightarrow{\overrightarrow{Y}} y_j$, $z_i \xrightarrow{\overrightarrow{Z}} z_j$ and $w_i \xrightarrow{\overrightarrow{W}} w_j$ denote the paths $y_i y_{i+1} \dots y_j$, $z_i z_{i+1} \dots z_j$ and $w_i w_{i+1} \dots w_j$, respectively. Throughout $h \xrightarrow{\overrightarrow{P}} g$ will denote a path starting from the vertex h and ending at the vertex g whose vertex set is to be specified. For $1 \leq i \leq s$, let $z_0 \xrightarrow{\overrightarrow{P}} z_i$ and $z_0 \xrightarrow{\overrightarrow{P}} w_i$ denote the paths both with vertex set $\{z_0, z_1, \dots, z_i, w_1, \dots, w_i\}$. Similarly, let $z_i \xrightarrow{\overrightarrow{P}} w_{s+1}$ and $w_i \xrightarrow{\overrightarrow{P}} w_{s+1}$ denote the paths both with vertex set $\{z_i, \dots, z_s, w_i, \dots, w_s, w_{s+1}\}$. Also, we let $z_0 \xrightarrow{\overrightarrow{P}} z_0 = z_0$ and $w_{s+1} \xrightarrow{\overrightarrow{P}} w_{s+1} = w_{s+1}$. Finally, given a sequence $p \xrightarrow{\overrightarrow{S}} q$, we denote by $q \xrightarrow{\overleftarrow{S}} p$ its reverse sequence. In the following Hamilton paths, strings like $a_i \xrightarrow{\overrightarrow{A}} a_j$ or $a_j \xrightarrow{\overleftarrow{A}} a_i$ with $i > j$ do not appear.

Case 2.1. (y_1, y_j) with $2 \leq j \leq k$.

$$y_1 \xrightarrow{\overrightarrow{Y}} y_{j-1} x z_1 z_0 w_1 w_2 \xrightarrow{\overrightarrow{P}} w_{s+1} y_k \xrightarrow{\overleftarrow{Y}} y_j.$$

Case 2.2. (y_i, y_j) with $2 \leq i < j \leq k$.

$$y_i \xrightarrow{\overrightarrow{Y}} y_{j-1} x y_{i-1} \xrightarrow{\overleftarrow{Y}} y_1 z_0 z_1 \xrightarrow{\overrightarrow{P}} w_{s+1} y_k \xrightarrow{\overleftarrow{Y}} y_j.$$

Case 2.3. (y_1, x) .

$$y_1 z_0 z_1 \xrightarrow{\overrightarrow{P}} w_{s+1} y_k \xrightarrow{\overleftarrow{Y}} y_2 x.$$

Case 2.4. (y_i, x) with $2 \leq i \leq k$.

$$y_i \xrightarrow{\overrightarrow{Y}} y_k w_{s+1} \xrightarrow{\overleftarrow{P}} z_1 z_0 y_1 \xrightarrow{\overleftarrow{Y}} y_{i-1} x.$$

Case 2.5. (y_i, z_j) with $1 \leq i \leq k-1$ and $0 \leq j \leq s$.

$$y_i \xrightarrow{\overleftarrow{Y}} y_1 x y_{i+1} \xrightarrow{\overrightarrow{Y}} y_k w_{s+1} \xrightarrow{\overleftarrow{P}} w_{j+1} w_j \xrightarrow{\overleftarrow{W}} w_1 z_0 \xrightarrow{\overrightarrow{Z}} z_j.$$

Case 2.6. (y_k, z_0) .

$$y_k \overleftarrow{Y} y_1 x z_1 \overrightarrow{Z} z_s w_{s+1} \overleftarrow{W} w_1 z_0.$$

Case 2.7. (y_k, z_j) with $1 \leq j \leq s$.

$$y_k \overleftarrow{Y} y_2 x y_1 z_0 \overrightarrow{P} w_{j-1} w_j \overrightarrow{W} w_{s+1} z_s \overleftarrow{Z} z_j,$$

where when $j = 1$ the string $z_0 \overrightarrow{P} w_{j-1}$ means z_0 .

Case 2.8. (y_1, w_1) .

$$y_1 z_0 z_1 x y_2 \overrightarrow{Y} y_k w_{s+1} \overleftarrow{P} w_2 w_1.$$

Case 2.9. (y_1, w_j) with $2 \leq j \leq s+1$.

$$y_1 z_0 w_1 \overrightarrow{W} w_{j-1} z_{j-1} \overleftarrow{Z} z_1 x y_2 \overrightarrow{Y} y_k w_{s+1} \overleftarrow{P} w_j.$$

Case 2.10. (y_i, w_j) with $2 \leq i \leq k$ and $1 \leq j \leq s+1$.

$$y_i \overrightarrow{Y} y_k x y_{i-1} \overleftarrow{Y} y_1 z_0 \overrightarrow{P} z_{j-1} z_j \overrightarrow{Z} z_s w_{s+1} \overleftarrow{W} w_j.$$

Case 2.11. (x, z_j) with $0 \leq j \leq s$.

$$x y_1 \overrightarrow{Y} y_k w_{s+1} \overleftarrow{P} w_{j+1} w_j \overrightarrow{W} w_1 z_0 \overrightarrow{Z} z_j.$$

Case 2.12. (x, w_j) with $1 \leq j \leq s+1$.

$$x y_k \overleftarrow{Y} y_1 z_0 \overrightarrow{P} z_{j-1} z_j \overrightarrow{Z} z_s w_{s+1} \overleftarrow{W} w_j.$$

Case 2.13. (z_i, z_j) with $0 \leq i < j \leq s$.

$$z_i \overleftarrow{P} z_0 y_1 x y_2 \overrightarrow{Y} y_k w_{s+1} \overleftarrow{P} w_{j+1} w_j \overrightarrow{W} w_{i+1} z_{i+1} \overrightarrow{Z} z_j.$$

Case 2.14. (z_0, w_1) .

$$z_0 z_1 x y_1 \overrightarrow{Y} y_k w_{s+1} \overleftarrow{P} w_2 w_1.$$

Case 2.15. (z_0, w_j) with $2 \leq j \leq s+1$.

$$z_0 w_1 \overrightarrow{W} w_{j-1} z_{j-1} \overleftarrow{Z} z_1 x y_1 \overrightarrow{Y} y_k w_{s+1} \overleftarrow{P} w_j,$$

Case 2.16. (z_i, w_j) with $1 \leq i < j \leq s+1$.

$$z_i \overrightarrow{Z} z_{j-1} w_{j-1} \overleftarrow{W} w_i w_{i-1} \overleftarrow{P} z_0 y_1 x y_2 \overrightarrow{Y} y_k w_{s+1} \overleftarrow{P} w_j,$$

where when $i = 1$ the string $w_{i-1} \overleftarrow{P} z_0$ means z_0 .

Case 2.17. (z_i, w_j) with $1 \leq j \leq i \leq s$.

$$z_i \overleftarrow{Z} z_j z_{j-1} \overleftarrow{P} z_0 y_1 x y_2 \overrightarrow{Y} y_k w_{s+1} \overleftarrow{P} w_{i+1} w_i \overleftarrow{W} w_j.$$

Case 2.18. (w_i, w_j) with $1 \leq i < j \leq s+1$.

$$w_i \overrightarrow{W} w_{j-1} z_{j-1} \overleftarrow{Z} z_i z_{i-1} \overleftarrow{P} z_0 y_1 x y_2 \overrightarrow{Y} y_k w_{s+1} \overleftarrow{P} w_j.$$

Thus we have shown that $H(n, \Delta)$ is hamiltonian-connected. Recall that any hamiltonian-connected graph of order at least 4 is 3-connected and hence has minimum degree at least 3. Since the graph $H(n, \Delta) - xz_1$ has connectivity 2, it is not hamiltonian-connected. For every edge $e \in E(H(n, \Delta))$ with $e \neq xz_1$, e has one endpoint of degree 3. Therefore $H(n, \Delta) - e$ has a vertex of degree 2, implying that it is not hamiltonian-connected. This completes the proof that $H(n, \Delta)$ is minimally hamiltonian-connected, and the theorem is proved. \square

Remark. The graphs $G(n, \Delta)$ and $H(n, \Delta)$ constructed in the above proof of Theorem 1 have degree sequences $\Delta, 3, 3, \dots, 3$ and $\Delta, 4, 3, \dots, 3$ respectively, and hence they have the minimum possible sizes among all graphs of order n with maximum degree Δ and minimum degree at least 3 in the two cases when $n - \Delta$ is odd and when $n - \Delta$ is even respectively. The graph constructed in [2] for Δ in the range $\lceil n/2 \rceil \leq \Delta \leq n - 3$ has degree sequence $\Delta, n - \Delta, 3, \dots, 3$.

Finally we pose two unsolved problems.

Problem 1. Let $n \geq 4$ be a given integer. What are the possible values of the minimum degree of a minimally hamiltonian-connected graph of order n ?

A computer search shows that every minimally hamiltonian-connected graph of order n with $4 \leq n \leq 10$ has minimum degree 3. The author does not know of an example of a minimally hamiltonian-connected graph with minimum degree at least 4. The following easier problem is of a more basic nature.

Problem 2. Does there exist a minimally hamiltonian-connected graph with minimum degree at least 4?

There are some sufficient conditions for hamiltonian-connected graphs; for recent ones see [1,3] and [4]. But very little is known about necessary conditions. Restrictions on the maximum or minimum degree of a minimally hamiltonian-connected graph may be viewed as necessary conditions for this smaller class of graphs.

Declaration of competing interest

There is no conflict of interest in this work.

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